# Data Structures and Algorithms COMP-251 A

## **Problem Assignment #4**

### 1. Balanced Search Trees

Consider the following sequence of keys: (5, 16, 22, 45, 2, 10, 18, 30, 50, 12, 1).

(a) Consider the insertion of items with this set of keys (in the order given) into an initiallyempty (2,4) tree T. Draw T after each insertion.

(b) Show that a (2,4) tree can have a different structure when the order of the items inserted is changed. Use the keys used in (a) in a different order of insertion.

(c) Prove that the height of a (2,4) tree storing *n* items is  $\Theta(\log n)$ , i.e., prove that is it both  $O(\log n)$  and  $\Omega(\log n)$ .

#### 2. Element Uniqueness (Problem 6.21 in Udi Manber text)

The input is a set *S* of *n* real numbers. Design an O(n) worst-case time algorithm to construct a number that is guaranteed **not** to be in *S*. Prove that  $\Omega(n)$  is a lower bound on the worst-case complexity of this problem under the *comparison-based* model of computation.

## **3.** Non-collinearity of Points

The input is a set *S* of *n* points in the plane in general position, i.e., no three points lie on a line. Invent an  $O(n \log n)$  worst-case time algorithm to construct a new point *X* such that *X* is not collinear with any pair of points in *S*. Prove the **correctness** of your algorithm. Prove the **complexity** of your algorithm.

#### 4. Maximal Points of a Set (Problem 8.18(a) in Udi Manber text)

The input is a set *S* of *n* points in the plane given by their *x* and *y* coordinates (real numbers). A point *p* in the plane is said to **dominate** another point *q* if **both** the *x* and *y* coordinates of *p* are greater than or equal to those of *q*. A point *p* is a **maximal** point of *S* if no other point of *S* dominates it. Invent an  $O(n \log n)$  worst-case time algorithm for finding **all** the maximal points of a given set *S*. Prove the **correctness** of your algorithm. Prove the **complexity** of your algorithm.